# The Grigorchuk and Grigorchuk-Machi Groups of Intermediate Growth 

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$$
\text { July 20, } 2019
$$

References:

1. de la Harpe, Topics in Geometric Group Theory, Chapter VIII.
2. Grigorchuk, Machi, "A group of intermediate growth acting by homomorphisms on the real line."
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## The Free Monoid

Definition
Let $A$ be a set. Then the free monoid on $A$, denoted $A^{*}$, is the set of all words over the alphabet $A$. More formally, the set of all finite sequences of elements of $A$.

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## Definition

The word length of an element $g \in G$ with respect to $S$ is

$$
|g|_{s}=\min \left\{|w| \mid w={ }_{G} g\right\} .
$$

## Growth of Groups

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\gamma_{G, S}(n)=\left|\left\{\left.g \in G| | g\right|_{S} \leq n\right\}\right|
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Exercise: $\gamma_{G \times H} \sim \gamma_{G} \gamma_{H}$.

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Exercise: $\gamma_{G \times H} \sim \gamma_{G} \gamma_{H}$.
Lemma (VIII.61,63)
If $\gamma_{G} \sim \gamma_{G}^{2}$, then there exists an $\alpha \in(0,1)$ such that $e^{n^{\alpha}} \preceq \gamma_{G}$.

## The Infinite Rooted Binary Tree

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Let $\mathrm{St}_{n}$ be the pointwise stabilizer of $T_{n}$. Note:

- $\operatorname{Aut}(T) / \mathrm{St}_{\mathrm{n}} \cong \operatorname{Aut}\left(T_{n}\right)$, so $\left[\mathrm{St}_{\mathrm{n}}: \operatorname{Aut}(T)\right]<\infty$.


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Therefore $\operatorname{Aut}(T)$ is residually finite.

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Similarly,

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\begin{aligned}
& \psi_{n}: \operatorname{Aut}(T) \rightarrow \operatorname{Aut}(T)^{2^{n}} \\
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Therefore $\operatorname{Aut}(T) \sim \operatorname{Aut}(T)^{2}$.
But $\operatorname{Aut}(T)$ is uncountable!

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$\begin{array}{llll}\square & ■ & ■ & ■ \\ ■ & ■ & \square & \square\end{array}$
$\because \quad . \quad!$
Now let $\Gamma=\langle a, b, c, d\rangle$.

## 「 via a Deterministic Finite Automaton



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$$
b(11)=1 c(1)=11 d(\varepsilon)=11
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$$
\begin{aligned}
d(110101) & =1 b(10101) \\
& =11 c(0101) \\
& =110 a(101) \\
& =1100 \mathrm{id}(01) \\
& =11000 \mathrm{id}(1) \\
& =110001 \mathrm{id}(\varepsilon) \\
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By recursive definition, $\psi_{1}: \mathrm{St}_{1, \Gamma} \rightarrow \Gamma^{2}$. Let $\psi_{1}=\left(\varphi_{0}, \varphi_{1}\right)$. Can check that

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\varphi_{0}= \begin{cases}b & \mapsto a \\ c & \mapsto a \\ d & \mapsto 1 \\ b^{a} & \mapsto c \\ c^{a} & \mapsto d \\ d^{a} & \mapsto b\end{cases}
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In particular, $\varphi_{0}: \mathrm{St}_{1} \rightarrow \Gamma$ is surjective. Therefore $\Gamma$ is infinite.

Recall $\psi_{1} \mid \mathrm{St}_{1}$ is injective.
Lemma (VIII.28)
$\left[\psi_{1}\left(\mathrm{St}_{1, \Gamma}\right): \Gamma^{2}\right]=8$.

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$\left[\psi_{1}\left(\mathrm{St}_{1, \Gamma}\right): \Gamma^{2}\right]=8$.
Therefore $\Gamma$ is commensurate to its square, and thus there exists $\alpha \in(0,1)$ such that $e^{n^{\alpha}} \preceq \gamma_{\Gamma}$.

## The Contracting Property

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## Lemma (VIII.62)

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Therefore for some $0<\alpha<\beta<1$,

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e^{n^{\alpha}} \preceq \gamma_{\Gamma} \preceq e^{n^{\beta}}
$$

so $\Gamma$ is a group of intermediate growth.

## Other Properties of $\Gamma$

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Theorem (Bartholdi '98, $\preceq_{2} \mid$ Erschler, Zheng '18, $\preceq_{1}$ )
Let $\lambda$ be the positive root of $x^{3}-x^{2}-2 x-4, \lambda \approx .7674 \ldots$
For all $\varepsilon>0$,

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－$\langle b, c, d\rangle \cong \mathbb{Z}_{2} \times \mathbb{Z}_{2}$ ．$\Gamma$ is 3－generated，but not 2－generated．
－$\Gamma$ is not finitely presentable．
－「 is torsion．Not obvious：Aut $(T)$ has elements of infinite order（exercise）．So，
－「 is a 2 －group．
－$\Gamma$ is not orderable．
－「 is not bounded torsion（ $\Leftarrow$ Zelmanov：bounded torsion +rf $\Rightarrow$ finite）．
－$\Gamma$ is amenable $(\Leftarrow$ subexp growth $)$ ．

## The Grigorchuk-Machi Group

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The proof of left-orderability uses Cantor's Theorem.
Theorem
If $X$ is a countable set with a total order $\leq x$ such that $\leq_{x}$ is dense and contains no first or last element, then $(X, \leq)$ is order-isomorphic to $(\mathbb{Q}, \leq)$.

Proof.
Exercise.

## The Grigorchuk-Machi Group

$$
\text { Let } \tilde{\Gamma}=\left\langle a, b, c, d \mid\left[a^{2}, b\right],\left[a^{2}, c\right],\left[a^{2}, d\right],[b, c],[b, d],[c, d]\right\rangle \text {. }
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Theorem (Grigorchuk '84)
The Grigorchuk-Machi group is of intermediate growth.

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## Corollary

Every left order on $\Gamma$ is Conradian (and one exists).

## Construction of $Q$

Start with
$\tilde{Q}=\left\langle a, b, c, d, \xi_{1}, \xi_{2}, \ldots\right|($ relations of $\tilde{\Gamma}),\left[a, \xi_{i}\right]$ for all $\left.i\right\rangle$.

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Define $R_{n}$ similarly to $K_{n}$, then set $R=\bigcup_{n=1}^{\infty} R_{n}$ and $Q=\tilde{Q} / R$ and $P=\tilde{P} / R$. We have that $\tilde{Q} / \tilde{P} \cong Q / P \cong\langle a\rangle \cong \mathbb{Z}$.

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Then prove that the order is dense and left-invariant.

Thank You!

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If you are from the U.S., happy Independence Day! Otherwise, happy 4th of July!

