

# The Sphynx kernel

```
;; Clock  
Computing  
<TnPr <TnPr  
End of computing.  
  
;; Clock -> 2002-01-17, 19h 25m 36s.  
Computing the boundary of the generator 19 (dimension 7) :  
<TnPr <TnPr <TnPr S3 <<Abar[2 S1][2 S1]>>> <<Abar>>> <<Abar>>>  
End of computing.  
  
Homology in dimension 6 :  
  
Component Z/12Z  
  
---done---  
  
;; Clock -> 2002-01-17, 19h 27m 15s
```

*Francis Sergeraert, Univ. of Grenoble and Lorient  
Logroño, May 2023*

## Semantics of colours:

**Blue** = “Standard” Mathematics

**Red** = Constructive, effective,  
algorithm, machine object, ...

**Violet** = Problem, difficulty,  
obstacle, disadvantage, ...

**Green** = Solution, essential point,  
mathematicians, ...

## Plan.

- 1. Introduction.
- 2. The general **CLOS** environment.
- 3. The toy example of implementing **magmas**  
through the **sphynx-class /MAGMA/**.
- 4. How it works.

# 1. Introduction.

Goal of the **Sphynx program**:

Designing a **computer algebra system**

convenient to **program computations**

in **algebraic topology**,

more precisely in **constructive algebraic topology**.

Must satisfy: **Code** production is **natural** and **easy**,

**Code modification and extension** are **natural** and **easy**,

**Code** is **easily readable**,

**Complete, safe and efficient typing**,

**Efficient solution for functional programming**,

**Easy switch efficient development / efficient production.**

Must be able to define

rich sets of complex mathematical classes,

with a complex hierarchy,

taking account of this hierarchy

to optimize the structure of this program

and also its use.

Warning: Using this kernel shows the goal

is essentially satisfied

But the kernel itself is quite complex.

The most important: **Safe** and **Efficient** Typing.

**Safe** = **Strict** and **arbitrary sophisticated typing**  
of the **elementary objects**.

Example:  $x \in K(\mathbb{Z}, 2)_n \Leftrightarrow x$  is a **list of integer lists**  
of lengths  $(n - 1, n - 2, \dots, 0)$ .

$$((7 \ 6 \ -6) \ (0 \ 5) \ (-10) \ ()) \in K(\mathbb{Z}, 2)_4$$

$$((7 \ 6 \ -6 \ ) \ (0 \ 5 \ 1) \ (-10) \ ()) \notin K(\mathbb{Z}, 2)_4$$

Safe = During **development**

the **type** of any **low level object** is **checked**,  
regardless of its **sophisticated type**.

Efficient = When the **development** is **finished**,

when the **program** is used  
for **actual computations**,  
these **type checkings** are **no longer used**.

## Brief “history” of typing:

Fortran-1970:  $i \dots, j \dots, \dots, n \dots = \text{integers}$   
other identifiers = float numbers

Lisp-1.5 (McCarthy): No typing at all.

PL/I, Pascal, C, C++, Caml, Common Lisp ...: Various weak typings.

Coq, HoTT, Lean, ...: Dependent types and strict typing.  
Programs completely proved.



## Brief “history” of typing:

Fortran-1970:  $i \dots, j \dots, \dots, n \dots = \text{integers}$

other identifiers = float numbers

Lisp-1.5 (McCarthy): No typing at all.

PL/I, Pascal, C, C++, Caml, Common Lisp ...: Various weak typings.

Sphinx kernel: Dependent types, strict typing and efficiency.

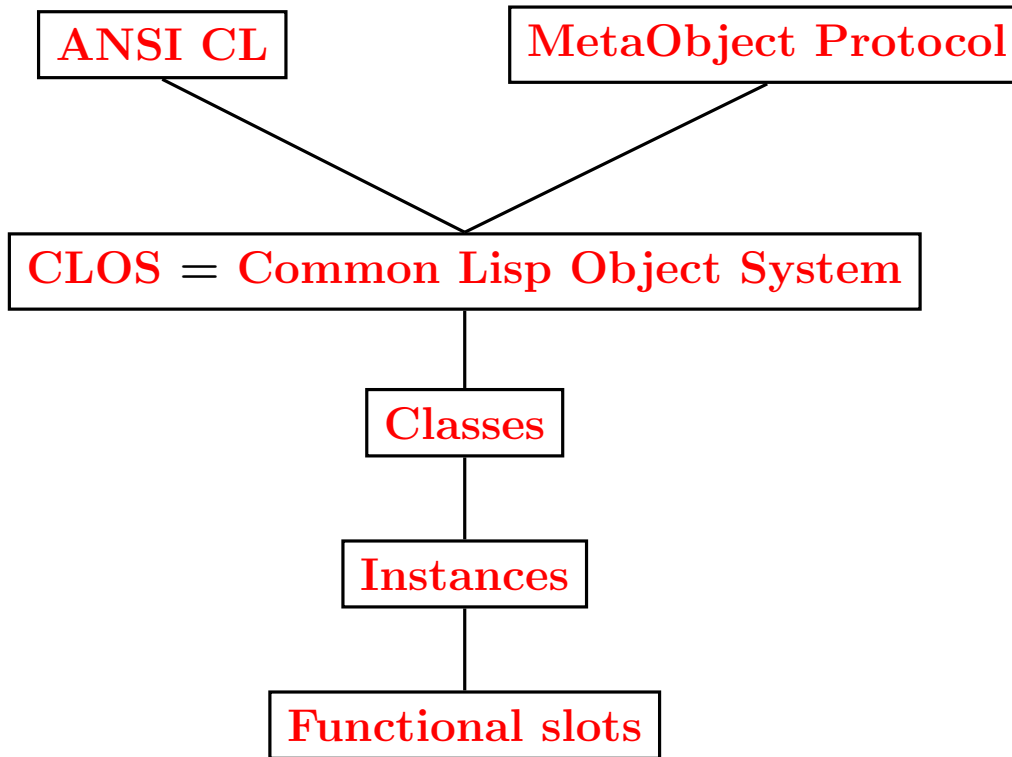
Programs reasonably reliable, but not completely proved.

Coq, HoTT, Lean, ...: Dependent types and strict typing.

Programs completely proved.

Efficiency?

## 2. The general CLOS environment.



## Rough history of CL OOP:

CL-1984 = First Steele's CL definition:

CLTL-1 (No OOP at all).

First OOP systems added to CL:

LOOPS, FLAVORS, OBJECT LISP, ...

1984-1990: MetaObject Protocol (Kiczales, Sussman, ...)

⇒ CLOS = Common Lisp Object System

CL-1990 = Final Steele's CL Definition  $\simeq$  ANSI-CL:

CLOS  $\subset$  ANSI-CL

CLOS  $\subset$  MetaObject Protocol

~~MetaObject Protocol  $\subset$  ANSI-CL~~

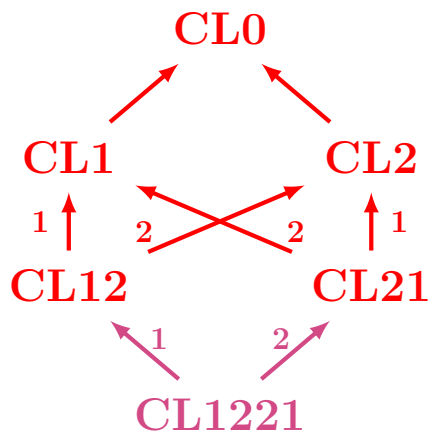
## Main components of CLOS:

- Classes.
  - **Multi-inheritance.**
  - Possible **redefinition.**
  - Elements = **Instances** = Collections of **slots.**
  - Powerful **initialization** of **slots.**
  - Complex hierarchies processed by **class-precedence-list.**
- Generic functions and methods.
  - Uses multi-inheritance.
  - Possible qualifiers **:before**, **:after** and **:around**
  - $\Rightarrow$  **Arbitrary complex** combinations  
of the **methods** of a **generic function.**

## Simplest diagrams of inheritance:

CL0 ← CL1

CL0 ← CL1 ← CL2



### 3. The toy example of Magmas.

Definition. A magma is a set

with a binary internal composition law  $*$ .

Example:  $M = \mathbb{N}$  with the composition law:

$$* : (x, y) \longmapsto x + 2y + xy$$

## Sphynx implementation of a set.

An **instance** of the **Sphynx class /SET/** has four essential **slots**:

1. **DFNT** (Definition): A **list** used only for documentation.
2. **MMBR** (Membership): A **predicate** strictly defining the **type** of the **elements**.
3. **CMPR** (Comparison): A **total order** for the **elements** of the **set**.
4. **CMPR=** : A **predicate** deciding the **equality** of two **elements**.

Example: the **set *NNI*** of the **non-negative integers**.

## Apparent Sphynx implementation of a Magma:

1. **DFNT** (as before).
2. **MMBR** (as before).
3. **CMPR** (as before).
4. **CMPR=** (as before).
5. **BNRL** : The **BiNaRy** composition **L**aw of this magma.

**MAGMA := SET + BNRL**



## Actual Sphynx implementation of a Magma:

1. **DFNT** (as before).
2. **MMBR** (as before).
3. **CMPR** (as before).
4. **CMPR=** (as before).
5. **BNRL: SOURCE:** (This-Magma This-Magma)  
**TARGET:** This-Magma  
**IMPL:** The actual Lisp functional object  
Other technical slots

**MAGMA := SET + BNRL**

#### 4. Technicalities of the SPHYNX Kernel.

SPHYNX Classes  $\in$  Graph under the top class /SPHYNX-OBJECT/.

Remark: Graph  $\neq$  Tree.

Any SPHYNX-CLASS is defined by SPHYNX-CLASS.

Any SPHYNX CLASS is a SubClass of /SPHYNX-OBJECT/.

A class hierarchy is defined by SPHYNX-CLASS and

computed by the MOP function:

COMPUTE-CLASS-PRECEDENCE-LIST

Defining a **new SPHINX-CLASS /CL/**.

Recipe:

- Use a **SPHYNX-CLASS** statement  
to define the **immediate super-classes**.  
and the **direct slots**.
- Use a **FUNCTIONAL-SLOT(S)** statement  
to define the possible functional slots.
- Use a **BIND-SLOTS** statement to define  
the **initialization process**  
of an **instance** of this **SPHYNX-CLASS**.

A **BIND-SLOTS** method can be **quite complex**.

How a `/SPHYNX-OBJECT/` of class `/CL/` is constructed:

- Allocate a fresh instance of class `/CL/`.
- Determine the class-precedence-list for `/CL/`.
- For every item of the class-precedence-list,  
call the corresponding `BIND-SLOTS` method.

Essential:

The order of `BIND-SLOTS` invocations is :

- From the most general classes
- to the most specific classes.

according to the class-precedence-list.

Implies in rare cases a **BIND-SLOTS** method  
must consider by anticipation  
a more specific class.

Typical example: the class of constructive sets = /SET/C.

Example: The set of the integers 3, 5 and 7.

Switching between

**DEVELOPMENT** and **PRODUCTION**

thanks to

the parameter **SYSDEF::\*CHECK-MEMBERSHIP-P\***.

To be used mainly for the **definitive compiled version**.

The END

```
;; Clock  
Computing  
<TnPr <TnPr  
End of computing.  
  
;; Clock -> 2002-01-17, 19h 25m 36s.  
Computing the boundary of the generator 19 (dimension 7) :  
<TnPr <TnPr <TnPr S3 <<Abar[2 S1][2 S1]>>> <<Abar>>> <<Abar>>>  
End of computing.  
  
Homology in dimension 6 :  
  
Component Z/12Z  
  
---done---  
  
;; Clock -> 2002-01-17, 19h 27m 15s
```

*Francis Sergeraert, Univ. of Grenoble and Lorient  
Logroño, May 2023*