

Sparse Permutations with Low Differential Uniformity¹

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Joint work with

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Outline

Preliminaries

Symmetric Cryptography
APN/AB Functions

Previous works

Permutations by Switching
Negative Answers

Main Results

Compositional Inverses
 $F_{s,1,\gamma}(x) = x^s + \gamma \text{Tr}(x)$
 $F(x) = x^{-1} + \gamma \text{Tr}(x^t)$

Conclusion

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Block Ciphers

$M \in \mathbb{F}_2^m$: plaintext,

$C \in \mathbb{F}_2^m$: ciphertext,

$K \in \mathbb{F}_2^k$: key.

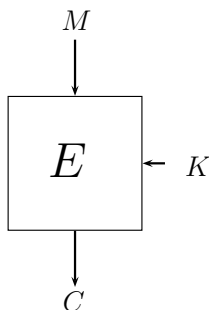
Block Cipher

$$E : \mathbb{F}_2^m \times \mathbb{F}_2^k \rightarrow \mathbb{F}_2^m$$

$$(M, K) \mapsto E(M, K) = C.$$

For a **fixed** key $K \in \mathbb{F}_2^k$,

$E_K(M) \mapsto C$, is a **permutation** of \mathbb{F}_2^m .



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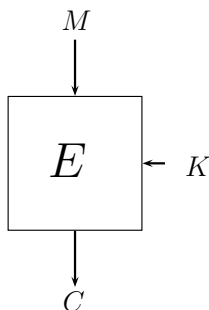
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Problems? In practice: $m \geq 64$ and $k \geq 80$ (!!!)

- ▶ For one key K , E_K permutes 2^{64} elements!

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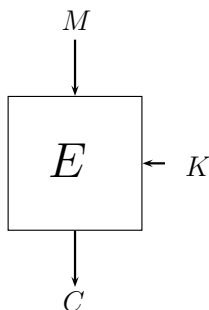
$C \in \mathbb{F}_2^m$: ciphertext,

$K \in \mathbb{F}_2^k$: key.

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- ▶ For one key K , E_K permutes 2^{64} elements!
- ▶ 2^{80} different permutations!

Substitution Permutation Networks

Add Round Key

$$\mathbb{F}_2^m \times \mathbb{F}_2^m \rightarrow \mathbb{F}_2^m$$

$$(M, \text{SubK}) \mapsto M \oplus \text{SubK}.$$

$$M = (M_0, \dots, M_{m/n-1}), M_i \in \mathbb{F}_2^n \text{ (word)}.$$

SBox (substitution)

Nonlinear Permutation:

$$S: \mathbb{F}_2^n \rightarrow \mathbb{F}_2^n$$

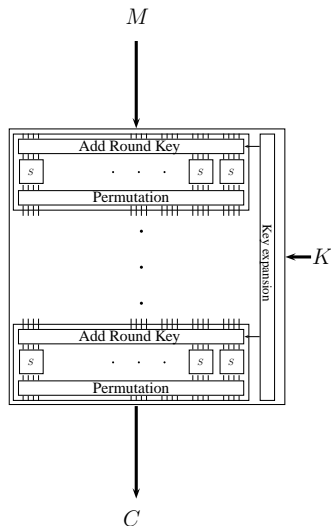
$$M_i \mapsto (S_0(M_i), \dots, S_{n-1}(M_i)).$$

In practice: $n = 4, 8$.

Permutation (diffusion)

Linear Permutation:

$$\mathbb{F}_2^m \rightarrow \mathbb{F}_2^m.$$



Unique Univariate Polynomial Representation of an SBox

$$\begin{aligned}
 F : \mathbb{F}_{2^n} &\rightarrow \mathbb{F}_{2^n} \\
 x &\mapsto \sum_{i=0}^{2^n-1} a_i x^i, \quad a_i \in \mathbb{F}_{2^n}.
 \end{aligned}$$

Definition

The **component functions** of a function F are

$$x \mapsto \text{Tr}(\lambda F(x)), \quad \lambda \in \mathbb{F}_{2^n}^*.$$

$\text{Tr} : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_2$ is the absolute trace.

Definition

The **algebraic degree** of a function F is

$$\deg(F) = \max\{\text{wt}(i) \mid a_i \neq 0\}.$$

$\text{wt} : \mathbb{F}_2^n \rightarrow \mathbb{N}$ is the binary Hamming weight.

F is said **sparse** if few of the a_i 's are nonzero.

Differential Uniformity

$$F : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$$

Definition

The **differential uniformity** of F is defined as

$$\delta(F) = \max_{a \neq 0, b \in \mathbb{F}_{2^n}} \#\{x \mid F(x) + F(x + a) = b\}.$$

F is called **APN** (almost perfect nonlinear) if $\delta(F) = 2$.

Facts about APN functions

APN functions have optimal resistance against *differential attacks*. Only **one APN permutation** is known for n even (Dillon's function, $n = 6$).

Nonlinearity

$$F : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$$

Definition

The **nonlinearity** of F is defined as

$$\mathcal{NL}(F) = 2^{n-1} - \max_{\alpha \neq 0, \beta \in \mathbb{F}_{2^n}} (2^{n-1} - wt(x \mapsto Tr(\alpha F(x) + \beta x))).$$

F is called **AB** (almost bent) if for all $\alpha \neq 0, \beta \in \mathbb{F}_{2^n}$

$$\mathcal{NL}(F) = 2^{n-1} - 2^{\frac{n-1}{2}}.$$

Facts about AB functions

Exist for n **odd** only!!

AB function have optimal resistance against *linear attacks*.

Any AB function is also APN.

Properties

$$F : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$$

Proposition

Differential uniformity and nonlinearity are invariants under:

- ▶ **compositional inversion** (when talking about permutations).
- ▶ **“Extended Affine”(EA)-equivalence** (i.e. $F \sim_{EA} G$ if there exist affine permutations A_0 and A_1 and some affine function A_2 such that $G = A_0 \circ F \circ A_1 + A_2$).

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Permutations by switching

P. Charpin and G. Kyureghyan [SETA08] [FFA09] [AMS10]:

Theorem

Let G be a permutation on \mathbb{F}_2 , f be any Boolean function on \mathbb{F}_2 .
Then

$$F(x) = G(x) + \gamma f(x), \quad \gamma \in \mathbb{F}_2^*,$$

is a permutation of \mathbb{F}_2^n **if and only if**

$$f \circ G^{-1}(x) + f \circ G^{-1}(x + \gamma) = 0 \quad \text{for all } x \in \mathbb{F}_2^n.$$

Let $\lambda \in \mathbb{F}_2^n$: $Tr(\lambda F(x)) = Tr(\lambda G(x)) + Tr(\lambda \gamma) f(x)$.

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Let $\lambda \in \mathbb{F}_2^n$: $\text{Tr}(\lambda F(x)) = \text{Tr}(\lambda G(x)) + \text{Tr}(\lambda \gamma) f(x)$.

Proposition

Let $\delta(G) = \rho$, then $\delta(F) \leq 2\rho$.

Particular case with monomials

Definition

Let $1 \leq s, t \leq 2^n - 2$ and $\gamma \in \mathbb{F}_{2^n}^*$.

$$F_{s,t,\gamma}(x) = x^s + \gamma \operatorname{Tr}(x^t).$$

Theorem

$F_{s,t,\gamma}$ is a **permutation** on \mathbb{F}_{2^n} if and only if

$$\gcd(s, 2^n - 1) = 1,$$

$t \equiv 2^j(2^i + 1)s \pmod{2^n - 1}$ for some $0 \leq i, j \leq n-1$, $i \neq n/2$,

and either **(a)** or **(b)** holds:

(a) $i = 0$ and $\operatorname{Tr}(\gamma) = 0$.

(b) $i > 0$ and $\gamma \in \mathbb{F}_{2^{\gcd(2i,n)}}$ with $\operatorname{Tr}(\gamma^{2^i+1}) = 0$.

Properties

$$F_{s,t,\gamma}(x) = x^s + \gamma \operatorname{Tr}(x^t).$$

Fact: $x \mapsto x^s$ is APN $\Rightarrow \gcd(s, 2^n - 1) = 3$ if n is even and 1 otherwise.

Proposition

- ▶ If n is *even* and $x \mapsto x^s$ is APN then F is not a permutation.
- ▶ If n is *odd* and $t = s(2^j + 1)$ where $\gcd(j, n) = 1$. Then F is not a permutation. It is a 2-to-1 function when $\gcd(s, 2^n - 1) = 1$ and $\gamma = 1$.

Proposition

There is **no permutation** on \mathbb{F}_{2^n} of the shape

$$F(x) = x^{2^j+1} + \gamma \operatorname{Tr}(x^{(2^j+1)(2^j+1)}),$$

with $\gcd(i, n) = 1$ and $\gcd(j, n) = 1$.

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$$F_{s,1,\gamma}(x) = x^s + \gamma \operatorname{Tr}(x)$$

$$F(x) = x^{-1} + \gamma \operatorname{Tr}(x^t)$$

Conclusion

Compositional inverses

$$F(x) = G(x) + \gamma f(x), \quad \gamma \in \mathbb{F}_{2^n}^*.$$

Theorem

If F is a **permutation** on \mathbb{F}_{2^n} (i.e. $f \circ G^{-1}(x) + f \circ G^{-1}(x + \gamma) = 0$), then

$$F^{-1} = G^{-1} \circ Q,$$

where $Q(x) = x + \gamma f \circ G^{-1}(x)$.

Because

$$F = Q \circ G \quad \text{and} \quad Q = Q^{-1}.$$

Compositional inverses

$$F_{s,t,\gamma}(x) = x^s + \gamma \operatorname{Tr}(x^t).$$

Theorem

Let $F_{s,t,\gamma}$ be a **permutation** on \mathbb{F}_{2^n} with $t = s(2^i + 1)$ for some i .
Let $\sigma = s^{-1} \pmod{2^n - 1}$. Then,

$$\begin{aligned} F_{s,t,\gamma}^{-1}(x) &= (x + \gamma \operatorname{Tr}(x^{2^i+1}))^\sigma \\ &= x^\sigma + \left(\sum_{j \prec \sigma} x^j \gamma^{\sigma-j} \right) \operatorname{Tr}(x^{2^i+1}). \end{aligned}$$

Definition: Let $a = \sum a_i 2^i$ and $b = \sum b_i 2^i$. If $a_i \geq b_i$ for all i then

$$b \preceq a \quad \text{and} \quad b \prec a \text{ if } a \neq b.$$

Special case when $t = 1$

$$F_{s,1,\gamma}(x) = x^s + \gamma \text{Tr}(x)$$

Theorem

$F_{s,1,\gamma}$ is a **permutation** on \mathbb{F}_{2^n} if and only if

1. $s = \frac{2^j}{2^i+1} \pmod{2^n - 1}^\dagger$ for some $i > 0$ and $j \geq 0$; **AND**
2. $\gcd(i, n) = \gcd(2i, n)$; **AND**
3. $\gamma \in \mathbb{F}_{2^{\gcd(i,n)}}$ such that $\text{Tr}(\gamma^{2^i+1}) = 0$.

Moreover (when $j = 0$),

$$F_{s,1,\gamma}^{-1}(x) = x^{2^i+1} + (\gamma^{2^i+1} + \gamma^{2^i}x + \gamma x^{2^i}) \text{Tr}(x^{2^i+1}).$$

In this case,

$$\delta(F_{s,1,\gamma}) = 2^{\gcd(i,n)}, \quad \mathcal{NL}(F_{s,1,\gamma}) = 2^{n-1} - 2^{(n+\gcd(i,n)-2)/2},$$

and $\deg(F_{s,1,\gamma}) = \frac{n-\gcd(i,n)+2^\dagger}{2}$, $\deg(F_{s,1,\gamma}^{-1}) = 3$.

† (On inversion in \mathbb{Z}_{2^n-1} [KS13])

Specific Classes

$$F_{s,1,\gamma}(x) = x^s + \gamma \text{Tr}(x)$$

Let $n = 2m$, with m odd. Take

1. $\gcd(i, n) = 2$;
2. $s = \frac{1}{2^i+1}$;
3. $\gamma \in \mathbb{F}_{2^2}^*$ such that $\text{Tr}(\gamma^{2^i+1}) = 0$ (i.e. $\gamma = 1$).

Then

$$F_{s,1,\gamma}(x) = x^{\frac{1}{2^i+1}} + \text{Tr}(x)$$

is a **permutation** on \mathbb{F}_{2^n} . Moreover

$$\delta(F_{s,1,\gamma}) = 4, \quad \mathcal{NL}(F_{s,1,\gamma}) = 2^{n-1} - 2^m \quad \text{and} \quad \deg(F_{s,1,\gamma}) = m.$$

Its *compositional inverse* is

$$F_{s,1,\gamma}^{-1}(x) = x^{2^i+1} + (1 + x + x^{2^i+1}) \text{Tr}(x^{2^i+1}).$$

With the multiplicative inverse function

Before getting started ...

Lemma

Let $F(x) = x^{-1} + \gamma f(x)$ where f is any boolean function on \mathbb{F}_{2^n} .

Then

$$\delta(F) \in \begin{cases} \{2, 4\} & \text{when } n \text{ is odd,} \\ \{4, 6\} & \text{when } n \text{ is even.} \end{cases}$$

It is worth noticing that

for n even, $\forall f$, $\forall \gamma$, F can not be APN.

Specific Classes

$$F(x) = x^{-1} + \gamma \text{Tr}(x^t)$$

Proposition

Let $1 \leq i \leq n$ with $i \neq n/2$, $\gamma \in \mathbb{F}_2^{\text{gcd}(2i,n)}$ such that $\text{Tr}(\gamma^{2^i+1}) = 0$, and

$$F_{i,\gamma}(x) = x^{-1} + \gamma \text{Tr}(x^{2^{n-1}-2^{i-1}-1}).$$

Then $F_{i,\gamma}(x)$ is a **permutation** on \mathbb{F}_{2^n} with

$$\delta(F_{i,\gamma}) \in \begin{cases} \{2, 4\} & \text{when } n \text{ is odd,} \\ \{4, 6\} & \text{when } n \text{ is even.} \end{cases}$$

Remark: when $\text{Tr}(\gamma^{2^i+1}) = 1$, $F_{i,\gamma}(x)$ is 2-to-1, with the same differential uniformity.

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Conclusion and Perspectives

- ▶ Little is known about the *nonlinearity* of the functions $F_{s,t,\gamma}$;
- ▶ Although we have a *good upper bound* on differential uniformity, it is generally *difficult* to get general result;
- ▶ Other classes of functions remain to be studied:
 - ▶ $x^{2^{2k}-2^k+1} + \gamma \text{Tr}(x^{2^{3i}+1})$,
 - ▶ $x^s + \gamma \text{Tr}(x^t + x^r)$,
 - ▶ \vdots

Thank you!